

On online Ramsey theory

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Ramsey Theory

Theorem (Ramsey's theorem, 1930)

For positive integers r and s , every graph of sufficiently large order has a complete graph of r vertices or an independent set of s vertices.

Theorem (Van der Waerden's theorem, 1927)

For positive integers c and n , every c -coloring of integers in $\{1, 2, \dots, N\}$ for sufficiently large N induces a monochromatic arithmetic progression of length n .

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An *online Ramsey game* for H on \mathcal{C} is a game between two players, Builder and Painter, with the following rules:

- ▶ Each turn:
 - ▶ Builder draws finitely many vertices and a new edge so that the resulting graph is in \mathcal{C}
 - ▶ Painter colors the new edge either red or blue.

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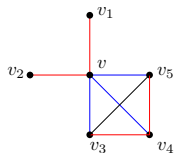
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- ▶ Each turn:
 - ▶ Builder draws finitely many vertices and a new edge so that the resulting graph is in \mathcal{C}
 - ▶ Painter colors the new edge either red or blue.
- ▶ If Builder can force Painter to make a monochromatic copy of H , then **Builder** wins.
- ▶ If Painter can avoid creating a monochromatic copy of H forever, then **Painter** wins.

Example

- ▶ **Builder** wins the online Ramsey game for C_3 on **planar graphs**.



Self-unavoidability

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Proposition (Grytczuk, Hałuszczak, and Kierstead, 2004)

Builder wins the online Ramsey game for every *k-colorable graph* on *k-colorable graphs*.

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However, this is not true for outerplanar graphs!

On Planar Graphs

Proposition (Grytczuk, Hałuszczak, and Kierstead, 2004)

Builder wins the online Ramsey game for C_n on *planar graphs* for all n .

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Conjecture (Grytczuk, Hałuszczak, and Kierstead, 2004)

Builder wins the online Ramsey game for H on *planar graphs* if and only if H is *outerplanar*.

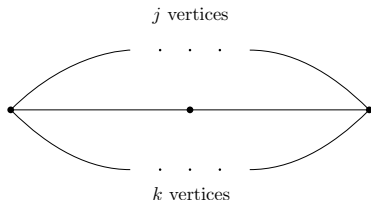
On Planar Graphs

Theorem (Petříčková, 2014)

For every *outerplanar* graph H , *Builder* wins the online Ramsey game for H on *planar graphs*.

Theorem (Petříčková, 2014)

Builder wins the online Ramsey game for $\theta_{2,j,k}$ on *planar graphs* for even j, k , while $\theta_{2,j,k}$ is *not outerplanar*.



Forbidden Minor Characterization

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Theorem (Grytczuk, Hałuszczak, and Kierstead, 2004)

Painter wins the online Ramsey game for C_3 on **outerplanar graphs**.

Theorem (Grytczuk, Hałuszczak, and Kierstead, 2004)

Builder wins the online Ramsey game for C_3 on **planar 2-degenerate graphs**.

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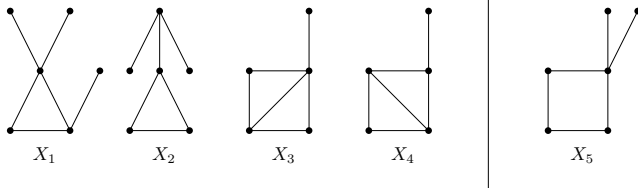
Theorem (C., Choi, Jeong, Oum, 2015+)

Painter wins the online Ramsey game for C_3 on **K_4 -minor-free graphs**.

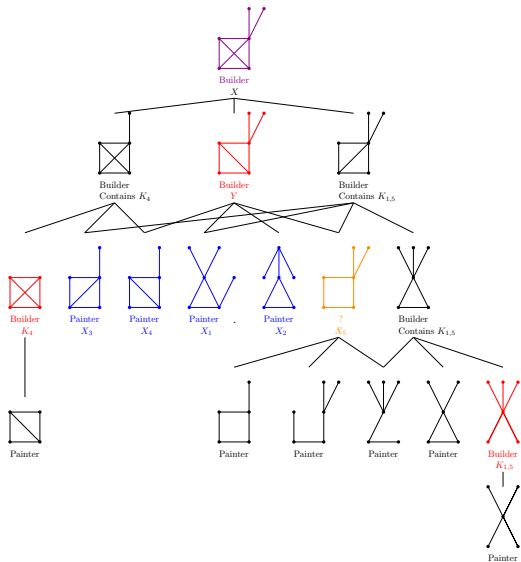
Forbidden Subgraph Characterization

Theorem (C., Choi, Jeong, Oum, 2015+)

Let F be a connected graph not isomorphic to X_5 . Painter wins the online Ramsey game for C_3 on F -free graphs if and only if F is isomorphic to a subgraph of X_i for some $1 \leq i \leq 4$.



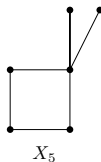
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Question

Who wins the online Ramsey game for C_3 on X_5 -free graphs?



Thank you very much.